Experiment on Cristian's and Berkeley Time Synchronization Algorithms





Centralized Algorithm:

1) Cristian's algorithm. (timeserver)

2) Berkeley algorithm. (coordinator)

Concept of Simulation Time

- Time changes on the basis of state change.
- Time always advances
- Time doesn't increases by the fix amount

Cristian's Algorithm

- External clock synchronization method.
- A process is the *time server* in the system.
- External time source *(Coordinated Universal Time)* is used as reference for synchronizing computer clocks with real time.
 - UTC is an international standard.
 - Standard bodies which, disseminate UTC signal by radio, telephone and satellite. One of the body is Geostationary Operational Environmental Satellites. (GEOS)

Cristian's Algorithm ... contd

- Other processes request for current time by sending request message (req_message) to time server.
- *time server* sends a reply by attaching current time with each reply message.
- Message suffers:
 - Transmission delay

Apparatus for Cristian's Algorithm :

- *time server* maintains a *message_queue* for arriving requests from other processes.
- A delay value is generated at *message_queue,* due to message queue delay (*states for which request was in queue*)
- Other issues : clock skew and clock drift
 - *clock skew :* The difference in time values of two clocks.
 - *clock drift :* The difference between speeds of two clocks.

Implementation for measuring performance....contd

- In *run()* step,
 - The messages are delivered in the increasing order of delay, to *time server*.
 - timeserver computes delay_at_rqst_queue (states for which the message was in queue)
 - *timeserver* sends a reply message to requesting process.
 - Requesting process receives the message and updates its time as follows:

current time = Tserver + (T1 - T0) / 2,

Tserver = server time returned by time server.

In simulation engine $T1 - T0 = message_queue_delay$

Berkeley Algorithm:

- coordinator polls other processes.
- Finds difference between its time and time of selected processes
- Then takes an average including its own time.
- Coordinator predicts time of other process with an error.
- Inform polled processes about correction (this message has error included)

Apparatus:

- coordinator polls processes and sends them request for there time
- processes sends reply.
- processes differ in their times
- Coordinator predicts error for each process.
- coordinator puts an upper bound errors .
- Two test environments are:
 - with minimal (or no) errors
 - with errors



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Future Work

- Enhance design for fault tolerance in distributed system.
- Measure the performance of the system by combining Cristian's and Berkley time synchronization algorithms.

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